

Profile and reorder code execution in Geant4 to increase performance

A Google Summer of Code Project

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- **Profile** Geant4 to identify potential targets of optimization (first half of GSoC period)
- **Reorder code execution** to improve serial performance (2nd half)

In reality Goals were interchangeable

- **Port of Geant4 to Solaris 11/amd64 to access DTrace** profiling tool
- **Tool to compare 2 versions** of an application and generate an HTML report
 - Tested on FullCMS and Simplified Calorimeter
 - Example (clickable): <http://island.quantumachine.net/~stathis/geant4/run-5550/smartstack.html>

- “D” stands for Dynamic- it dynamically instruments a running program, by modifying its instructions while it is executing
- **Deep inspection**
 - Arbitrary instructions
 - CPU registers
 - CPU hardware counters, etc
- **Sophisticated profiling** (e.g., speculative tracing)
- **Built-in aggregation** functions
 - count, sum, avg, min, max, stddev, {l,}quantize
- **Negligible runtime overhead**

- **Safe** to use in production environments
 - Safety was one of the central architectural decisions upon DTrace was built
 - Explains why some common language constructs aren't supported (e.g., for-loops)
- **No source code modification** of the profiled application needed
- Can operate on **multithreaded** programs (has support for thread-local variables)
- Runs on **Mac OSX** out of the box; Linux port is on the way
- Profiling done via a simple language called D (resembling C and awk)
 - Scripts can be shared, reviewed, reused, made be run unattended

Some of the ideas explored

- **Particle bunching (G4SmartTrackStack)**
- Hard-coded stepping manager (G4SteppingManager)
- Caching of cross-sections calculations in hadronic processes (G4CrossSectionDataStore)
- Reducing branch mispredictions in Value() (G4PhysicsVector)
- Caching values of $\ln(\text{Energy})$ (G4Track)

Definition Process *same* particle types before switching to another particle type. E.g.,

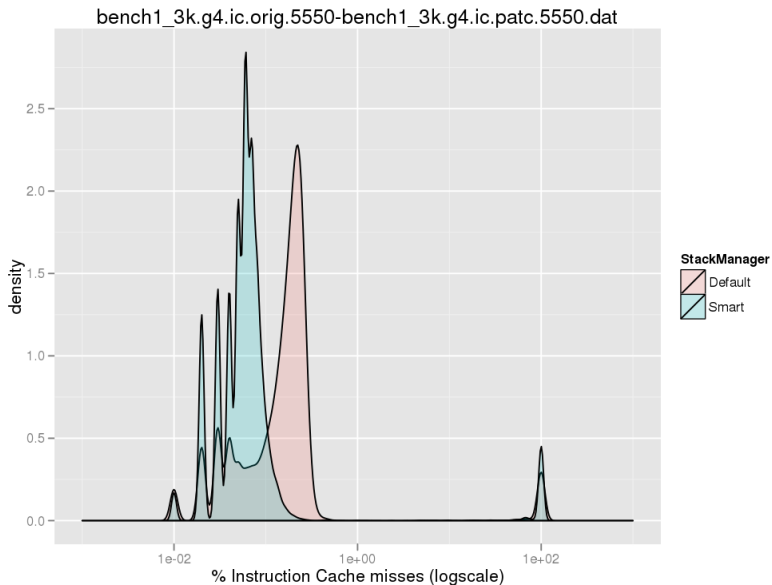
$$\dots, e^-, e^-, \dots, e^-, \gamma, \gamma, \dots, \gamma, \dots$$

Why Better *cache utilisation* due to access to the same physics list

Number of stacks we are using: 5

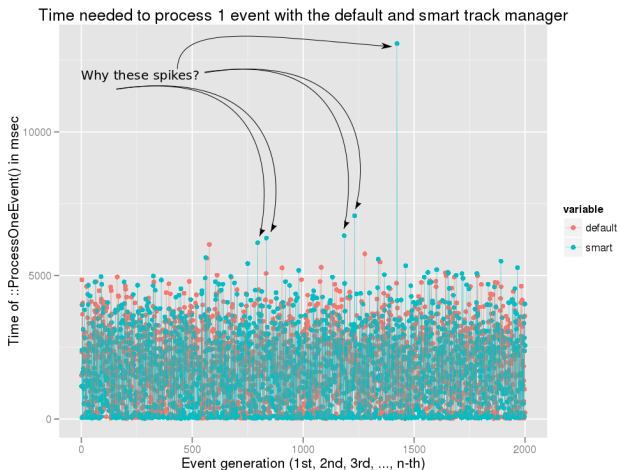
- 1 Primary particles + everything not belonging to:
- 2 Neutrons
- 3 Electrons
- 4 Gammas
- 5 Positrons

Particle "bunching" 2/2



Speculative tracing - A real use case

Problem Some ProcessOneEvent() need much more than average time to complete



Speculative tracing - A real use case

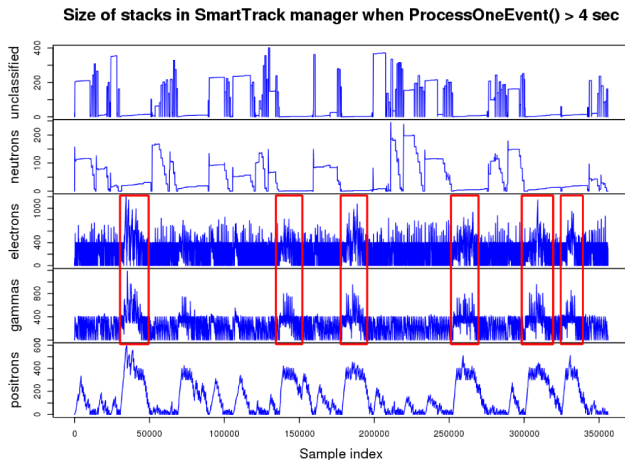
Strategy We are going to trace all `ProcessOneEvent()` calls, but commit to our tracing buffer *only* those that behave bad.

In this context, "trace" refers to looking at stacks' sizes when `ProcessOneEvent()` stalls while processing the event.

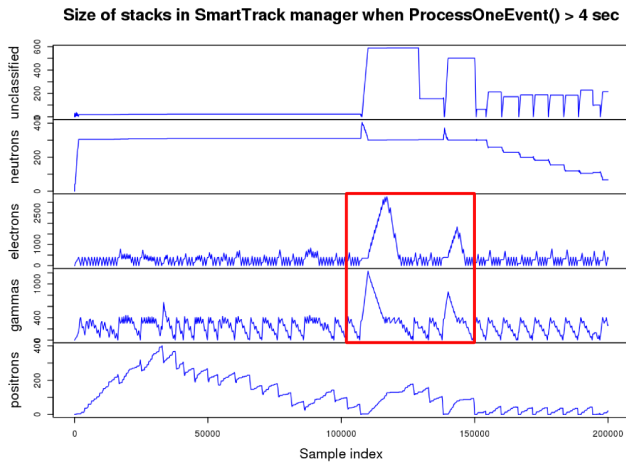
Speculative tracing - A real use case cont.

Hint The maximum desired size for all stacks was requested to be 400.

e^- and γ too often will not honour that limit.



Speculative tracing - A real use case cont. - Zoom 1/2



Problem A flamegraph showing CPU utilization identified cross-section calculations in hadronic processes as a significant contributor

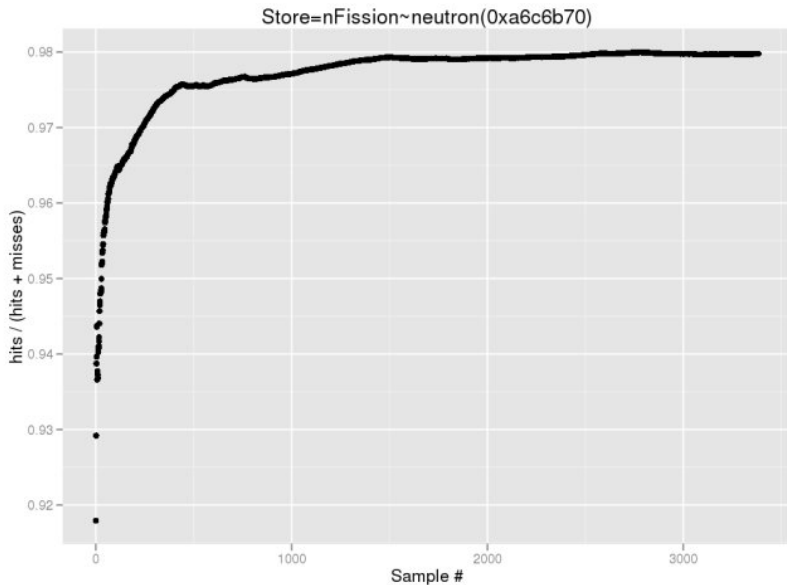
Idea Cache the values on some bin energy level

Result After many iterations, we have a version where the **hits ratio are very high** and there's **probably a benefit of a few percent** (not yet quantified)

TODO Run enough simulations to extract the benefit. Study the **ramifications of bin'ing the energy** from the physics POV.

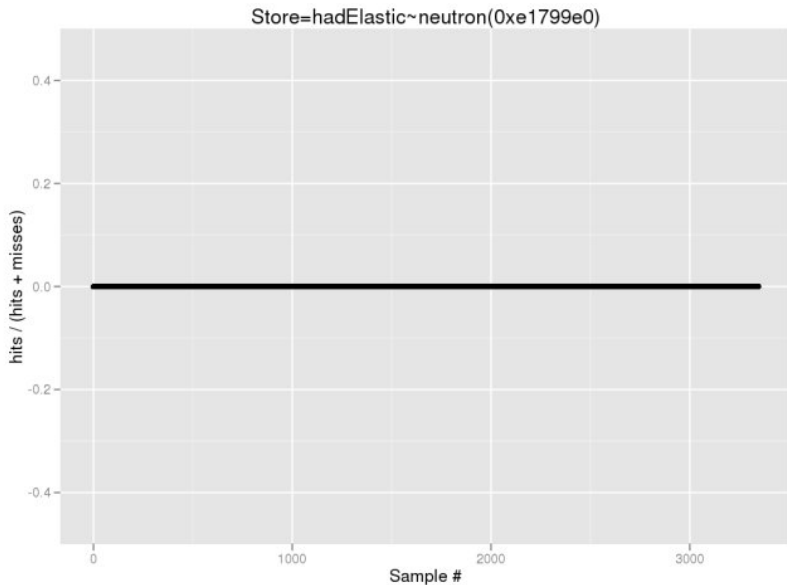
Not all hadronic processes are cache-friendly 1/2

<http://island.quantummachine.net/~stathis/geant4/hits>



Not all hadronic processes are cache-friendly 2/2

<http://island.quantummachine.net/~stathis/geant4/hits>



"Problem" A flamegraph showing branch mispredictions identified G4PhysicsVector::Value() as a significant offender

Idea Try to collapse some of the if-blocks, gaining branch predictability, but executing more cpu instructions

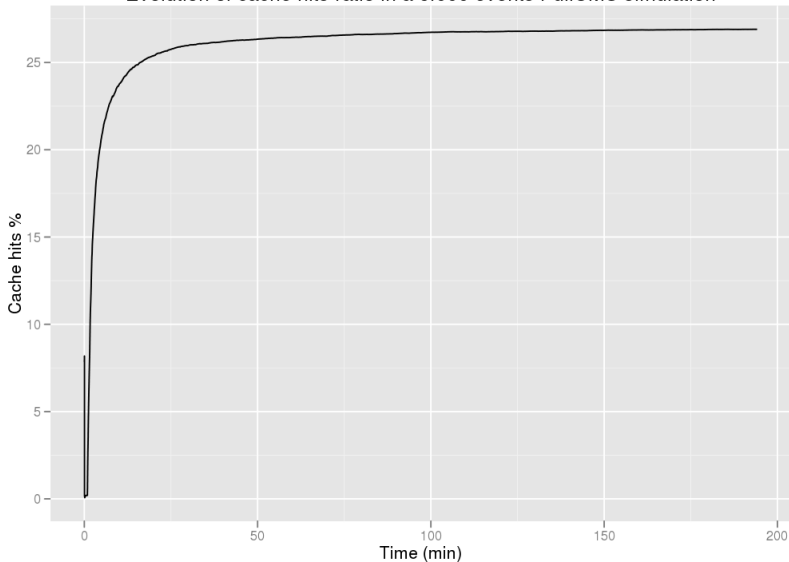
Result The branch mispredictions reduced (expected), but the average time spent in that function was actually larger

Objective Calculate the cache hits ratio in G4PhysicsVector::Value()

```
# dtrace -qn '  
/* 0xc0 is the offset inside Value() where a fast cache hit takes place */  
pid$target::_ZN15G4PhysicsVector5ValueEd:c0  
{  
    @branch = count()  
}  
  
pid$target::_ZN15G4PhysicsVector5ValueEd:entry  
{  
    @total = count()  
}  
  
tick-100ms  
{  
    printa(@branch)  
    printa(@total)  
}' -c '/home/stathis/geant4.9.5.p01/bin/full_cms ./bench1_5k.g4' -o val
```

Reducing branch mispredictions in G4PhysicsVector::Value()

Evolution of cache hits ratio in a 5.000 events FullCMS simulation



- The benefit of caching outweighs (as reality dictates) the penalty of branch mispredictions
- The eventual ratio is higher than that I had initially in mind
- **Lesson learnt:** let the system reach its equilibrium before drawing any conclusions
- **Lesson learnt:** if you optimize 1 micro-benchmark, you may hurt another (or more)

Enter the "rabbit" hole

- **Question** ::Value() has many distinct branches. How **fast** are compared to each other ?
- **Question** ::Value() has many distinct branches. How **many times** is each one executed ?

I will skip the DTrace script which is a bit long for a slide, but here are the graphs:

Reducing branch mispredictions in G4PhysicsVector::Value()

A preliminary analysis with DTrace showed that the anticipated benefit would be less than 1%, so this idea hasn't been actively pursued until now

Thank you. Questions?